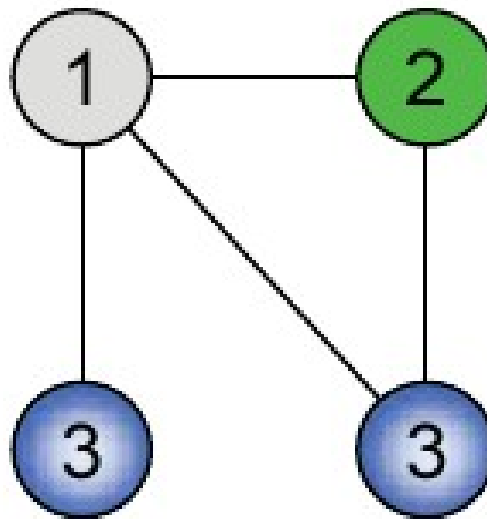


Vertex coloring

Problem

Problem 1.1 (Vertex Coloring). *Given an undirected graph $G = (V, E)$, assign a color c_u to each vertex $u \in V$ such that the following holds: $e = (v, w) \in E \Rightarrow c_v \neq c_w$.*



A non distributed solution

Assumption 1.2 (Node Identifiers). *Each node has a unique identifier, e.g., its IP address. We usually assume that each identifier consists of only $\log n$ bits if the system has n nodes.*

Algorithm 1 Greedy Sequential

- 1: **while** \exists uncolored vertex v **do**
 - 2: color v with the minimal color (number) that does not conflict with the already colored neighbors
 - 3: **end while**
-

Performance

Theorem 1.5 (Analysis of Algorithm 1). *The algorithm is correct and terminates in n “steps”. The algorithm uses $\Delta + 1$ colors.*

- Δ is the maximum degree of the graph
- The *chromatic number* $\chi(G)$ of G is the minimum number of colours in a proper vertex coloring of G

A useful subroutine

Procedure 2 First Free

Require: Node Coloring {e.g., node IDs as defined in Assumption 1.2}

Give v the smallest admissible color {i.e., the smallest node color not used by any neighbor}

- **Caveat:** no two nodes are coloured at the same time

Algorithm for synchronous case

Algorithm 3 Reduce

- 1: Assume that initially all nodes have ID's (Assumption 1.2)
 - 2: **Each node** v executes the following code
 - 3: node v sends its ID to all neighbors
 - 4: node v receives IDs of neighbors
 - 5: **while** node v has an uncolored neighbor with higher ID **do**
 - 6: node v sends “undecided” to all neighbors
 - 7: node v receives new decisions from neighbors
 - 8: **end while**
 - 9: node v chooses a free color using subroutine **First Free** (Procedure 2)
 - 10: node v informs all its neighbors about its choice
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- Thm.: algorithm is correct and has time complexity n . It uses $\Delta + 1$ colours

Trees

Lemma 1.9. $\chi(\text{Tree}) \leq 2$

Algorithm 4 Slow Tree Coloring

- 1: Color the root 0, root sends 0 to its children
 - 2: **Each node** v concurrently executes the following code:
 - 3: **if** node v receives a message x (from parent) **then**
 - 4: node v chooses color $c_v = 1 - x$
 - 5: node v sends c_v to its children (all neighbors except parent)
 - 6: **end if**
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- Caveat: how do we choose a root?

Trees/cont.

- The previous algorithm also works in the asynchronous case
 - Time complexity is the tree height
 - Message complexity is $n-1$
- Is it possible to do better?
- \log^*n time is possible!