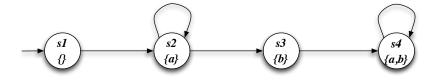


- 1. Express it in FOL.
- 2. Express it in *DL-Lite*<sub>A</sub>, highlighting parts that are not expressible.
- 3. Given the ABox  $A = \{C(c)\}$  and the boolean conjunctive query  $q(x) \leftarrow Rab(x, y), Rab(y, z), A(z)$ , return the certain answer by exploiting the *DL-Lite*<sub>A</sub> rewriting algorithm.

**Exercise 2.** Model check the Mu-Calculus formula  $\mu X.\nu Y.(a \vee [next]X) \wedge [next]Y$  and the CTL formula AFAGa against the following transition system:

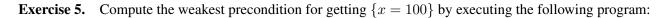


**Exercise 3.** Check whether CQ  $q_1$  is contained in CQ  $q_2$ , reporting canonical DBs and homomorphism:

$$\begin{array}{rcl} q_1(x_r) & \leftarrow & e(x_r, x_g), e(x_g, x_b), e(x_b, x_r). \\ q_2(x) & \leftarrow & e(x, y), e(y, z), e(z, x), e(z, v)e(v, w), e(w, z). \end{array}$$

**Exercise 4.** Compute the certain answers to the CQ  $q(x) \leftarrow M(x, y), E(y)$  over the incomplete database (naive tables):

E(mployee)	M(anager)	
name	mgr	mgd
Smith	Smith	$null_1$
$null_1$	$null_1$	Brown
Brown	Brown	$null_2$



```
x := y + 50;
if (y > 0) then
x := y + 100
else x := y + 200;
x := x + y;
```