SDCI 2012

HOT TOPICS IN SECURE AND DEPENDABLE COMPUTING FOR CRITICAL INFRASTRUCTURES JANUARY 15TH – 19TH, CORTINA D'AMPEZZO, ITALY

COLLABORATIVE EVENT PROCESSING FOR THE PROTECTION OF CRITICAL INFRASTRUCTURES TIME WINDOW BASED COMPUTATIONS: A BATCH APPROACH

Leonardo Aniello aniello@dis.uniroma1.it



DIPARTIMENTO DI INFORMATICA E SISTEMISTICA "A. RUBERTI"



CONTEXT: CYBER ATTACKS DETECTION

- Complex Event Processing (CEP)
 - enables the detection of complex attack patterns
 - huge amount of input data
 - timeliness requirements
- Time Window Based Computation (TWBC)
 - an attack takes place within a certain time window
 - we need a bound to the data to process

TIME WINDOW BASED COMPUTATION (TWBC)

- a particular elaboration of a set of events that happened within a specific time window
- events are fed in the form of an event stream
- focus on
 - windows with fixed-length T
 - a new window is started ΔT after the previous window was started

APPROACHES

• online

- events are processed as they enter the engine
- many CEP engines publicly available (System S, Esper, Drools Fusion)
- availability of mechanisms to properly manage TWBCs

• batch

- events are stored as they enter the engine
- stored events are processed periodically
- each computation concerns a specific time window
- MapReduce paradigm: Hadoop

HADOOP (1/2)

- distributed processing framework
- a computation is called job and its work is decomposed in map tasks and reduce tasks
- HDFS is the default storage, a distributed file system where data is organized in files and directories
- input data is fed to a variable set of map tasks
- they perform part of the overall elaboration and produce an halfway output
- such output is then computed by a fixed number of reduce tasks to produce the final output

HADOOP (2/2)

- limited slots for running map tasks concurrently
- if slots are exhausted, some tasks have to wait
 ⇒increase overall computation latency
- the number of required map tasks depends on how input data is organized in files
 - in general, the higher is the number of files, the higher is the number of map tasks
 - it's better a small number of large files than a large number of small files

Why Batch instead of Online

- fault tolerance
 - if an online engine crashes, the state is lost
 - if Hadoop crashes, events are still available on HDFS
- multiple sources
 - online engines usually have a single entry point: possible bottleneck
 - in Hadoop each source can feed data to the nearest Hadoop node
- availability of online services for MapReduce computations

TWBCs with Hadoop

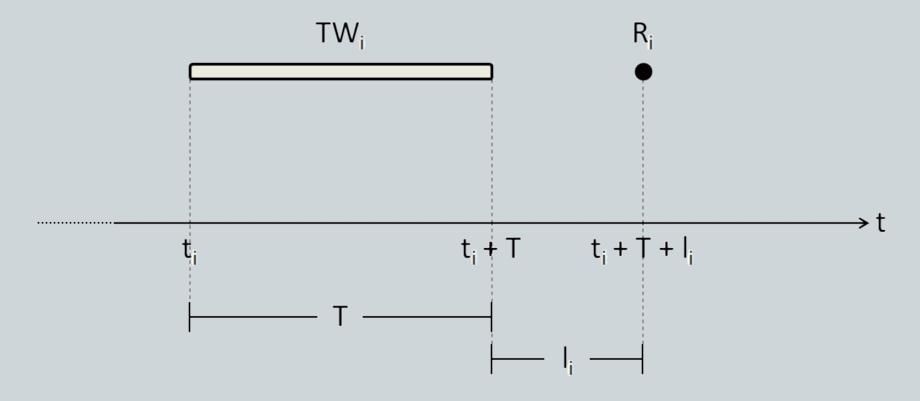
- we consider a single source
- events are written to an HDFS file as they enter the system
- once the file is closed, contained events are available for the elaboration
- rolling strategy

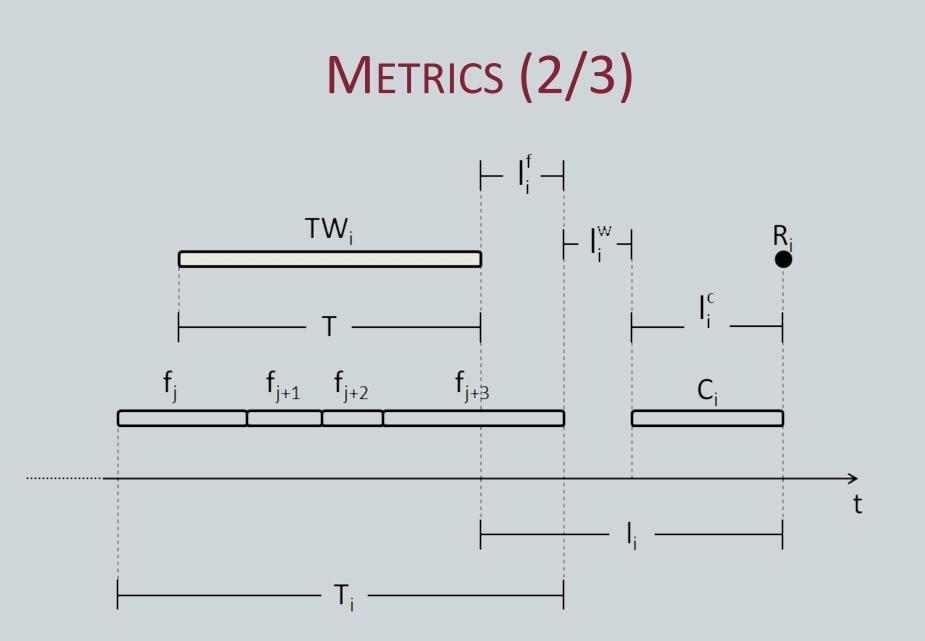
key role of the policy used to close HDFS files, it affects

- the number of files to process
- the delay the events are at disposal for the job

METRICS (1/3)

 how long does it take after the end of a time window to have the result of the corresponding computation?





METRICS (3/3)

• file latency

the time between the end of a time window and the closure of the file containing the last event of such window

• wait latency

the time between the closure of the file containing the last of event of a window and the begin of the correspondent computation

• computation latency

• time efficiency

the ratio T/T_i ; it expresses the percentage of data read from file that actually belong to TW_i

OPTIMIZE THE METRICS (1/2)

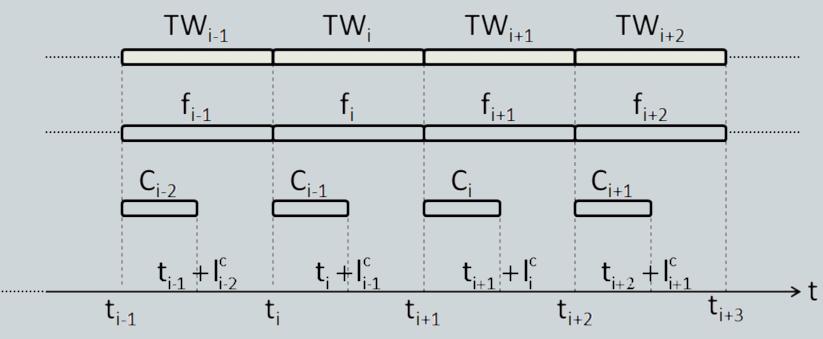
- file latency can be minimized by closing currently open file as soon as a time window ends
- wait latency can be minimized by starting a new computation on TW_i as soon as the file containing the last event of TW_i is closed
- time efficiency can be made as near as possible to 1 by
 - opening a new file just before the first event of a time window
 - minimizing file latency

OPTIMIZE THE METRICS (2/2)

- how can we reduce computation latency only acting on the way we manage HDFS files?
 - 1. decrease the number of files used to cover a TW
 - lower overhead due to map tasks management
 - lower probabilities to run out of map slots
 - 2. optimize time efficiency
 - since Hadoop is IO bound, reading only what is really required helps to decrease the latency

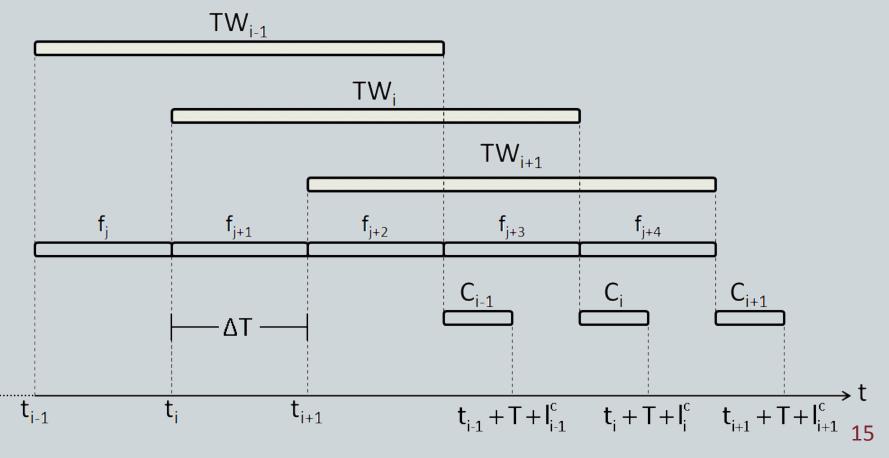
ROLLING STRATEGY FOR JUXTAPOSED TIME WINDOWS

- $\Delta T = T$; a single file for each time window
- file latency, wait latency and time efficiency optimized
- number of files per time window optimized



ROLLING STRATEGY FOR INTERLEAVED TIME WINDOWS

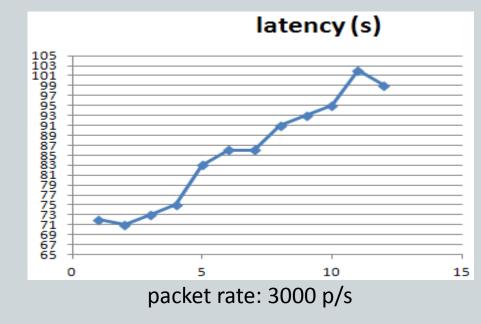
- $N \cdot \Delta T = T$; N files for each time window
- file latency, wait latency and time efficiency optimized



ROLLING STRATEGY FOR INTERLEAVED TIME WINDOWS EVALUATIONS

- T = 1 hour; Δ T = 1 minute; N = 60
- input data: fixed packet rate
- 7 nodes Hadoop cluster + 1 data source node
- analyze the latency varying the packet rate (12 time windows)

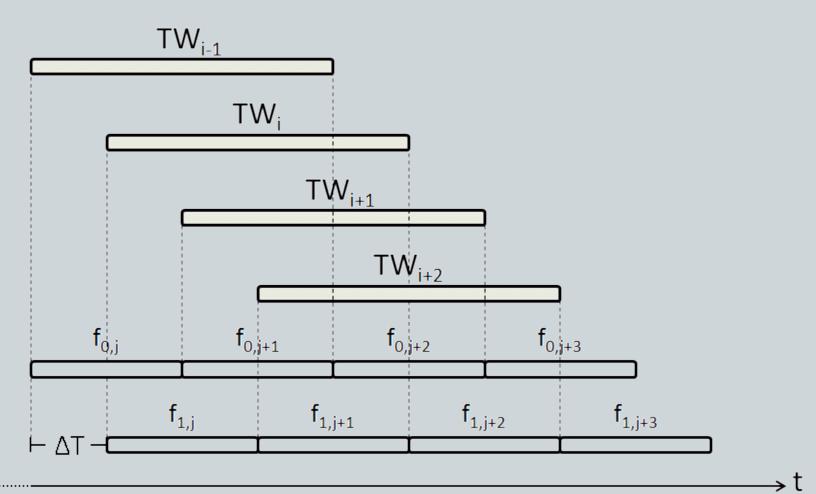
packet rate (p/s)	average latency (s)
500	43
1000	50
2000	63



MULTI FILES FLOWS ROLLING STRATEGY FOR INTERLEAVED TIME WINDOWS (1/4)

- to decrease the number of files per time window, we have to increase the time coverage of each file
- to keep files correctly fit time windows, so that both file latency and time efficiency are optimized, we have to use distinct files flows
- each files flow starts with an offset of ΔT from the previous one
- each time window has exactly one files flow which correctly fits

MULTI FILES FLOWS ROLLING STRATEGY FOR INTERLEAVED TIME WINDOWS (2/4)

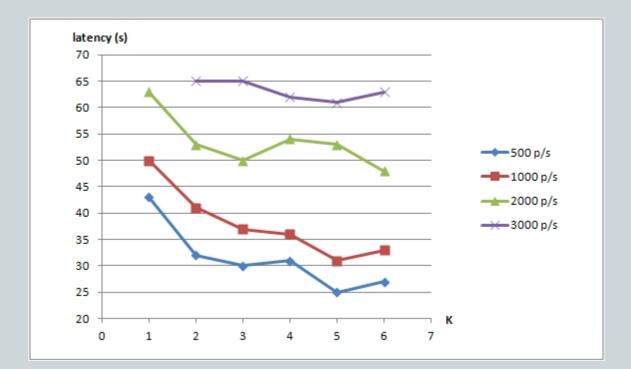


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MULTI FILES FLOWS ROLLING STRATEGY FOR INTERLEAVED TIME WINDOWS (3/4)

- K files flows (0 to K-1), where N mod K = 0
- file flow j includes files f_{j,0}, f_{j,1}, f_{j,2}, ...
- each file covers a K· Δ T time interval
- each time window is covered by N/K files
- we need more disk space (K times)

MULTI FILES FLOWS ROLLING STRATEGY FOR INTERLEAVED TIME WINDOWS (4/4)



average latency varying input packet rate and K

FUTURE WORKS

- go on with other evaluations
- consider multiple sources

Thank you!