Dynamic and Temporal Answer Set Programming on Linear Finite Traces

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The logical foundations of Answer Set Programming (ASP; (Lifschitz, 1999)) rest upon the logic of Here-and-There (HT; (Heyting, 1930)), or more precisely its equilibrium models (Pearce, 1997) that correspond to stable models semantics (Gelfond and Lifschitz, 1988). For defining extensions to ASP from firm logical principles, it has thus become good practise to first elaborate upon them in the setting of HT in order to afterwards consider the respective language fragments that are well suited in the context of logic programming.

This avenue was also followed in (Cabalar and Vega, 2007), which gave rise to the temporal extension of HT called Temporal Here-and-There and its non-monotonic counterpart Temporal Equilibrium Logic (for short THT and TEL (Aguado et al., 2013)). More precisely, TEL builds upon an extension of the logic of HT with Linear Temporal Logic (LTL; (Pnueli, 1977)). This results in an expressive non-monotonic modal logic, which extends traditional temporal logic programming approaches (Cabalar, Diéguez, and Vidal, 2015) to the general syntax of LTL and possesses a computational complexity beyond LTL (Bozelli and Pearce, 2015). As in LTL, a model in TEL is an infinite sequence of states, called a trace. However, this rules out computation by ASP technology (and necessitates model checking) and is unnatural for applications like planning, where plans amount to finite prefixes of one or more traces (cf. (Baier and McIlraith, 2006; De Giacomo and Vardi, 2013)).

Unlike this, we recently proposed in (Cabalar et al., 2018) an alternative combination of the logics of HT and LTL whose semantics rests upon finite traces. On the one hand, this amounts to a restriction of THT and TEL to finite traces. On the other hand, this is similar to the restriction of LTL to LTLf advocated by (De Giacomo and Vardi, 2013); see also (Baier and McIlraith, 2006). Our new approach, dubbed TELf, has the following advantages. First, it is readily implementable via ASP technology. Second, it can be reduced to a normal form which is close to logic programs and much simpler than the one obtained for TEL. Finally, its temporal models are finite and offer a one-to-one correspondence to plans. Interestingly, TELf also sheds light on concepts and methodology used in incremental ASP solving when understanding incremental parameters as time points.

Another distinctive feature of TELf is the inclusion of future as well as past temporal operators. We associate this with the following benefits. When using the causal reading of program rules, it is generally more natural to draw upon the past in rule bodies and to refer to the future in rule heads. A similar argument was put forward by Gabbay (1987) in his proposal of “declarative past and imperative future.” This format also yields a simpler normal form and lends itself to a systematic modeling methodology which favors the definition of states in terms of the past rather than mixing in future operators. For instance, in reasoning about actions, the idea is to derive action effects for the current state and check their preconditions in the previous one, rather than to represent this as a transition from the current to the next state. This methodology aligns state constraints, effect axioms, etc. to capture the present state. As well, past operators are much easier handled computationally than their future counterparts when it comes to incremental reasoning, since they refer to already computed knowledge.

TELf is implemented in the telingo system, extending the ASP system clingo to compute the temporal stable models of (non-ground) temporal logic programs. To this end, it extends the full-fledged input language of clingo with temporal operators and computes temporal models incrementally by multi-shot solving (Gebser et al., 2018) using a modular translation into ASP. telingo is freely available at github\(^1\). The interested reader might have a good time playing with the examples given in the examples folder at the same site.

Similar to the extension of LTLf to its (linear) dynamic logic counterpart LDLf (De Giacomo and Vardi, 2013), we just introduced in (Bossier et al., 2018) a dynamic extension of HT that draws upon this linear version of dynamic logic. We refer to the resulting logic as \((\text{Linear})\) Dynamic Logic of Here-and-There (DHT for short). As usual, the equilibrium models of DHT are used to define temporal stable models and induce the non-monotonic counterpart of DHT, referred to as \((\text{Linear})\) Dynamic Equilibrium Logic (DEL). In doing so, we actually parallel earlier work extending HT with LTL, ultimately leading to THT and TEL.

In fact, we show that THT (and its equilibrium counter-

\(^1\)https://github.com/potassco/telingo{}
part TEL) can be embedded into our new logic DHT (and DEL, respectively) — just as LTL can be put in LDL. Moreover, we prove that the satisfiability problem in DEL is EXPSpace-complete; it thus coincides with that of TEL but goes beyond that of LDL and LTL, both being PSPACE-complete. In fact, the membership part of this result is obtained by means of an automata-based method for computing DEL models. Finally, we show that the monotonic base logic of DEL, namely DHT, allows us to decide strong equivalence in DEL; this reinforces the adequacy of the relation between both logics.

These recent results open several interesting topics for future study. First, the version of DEL for finite traces, DEL_f, seems a natural step to follow, similar to the relation of LDL and LDf. We plan to propose and analyse this variation in an immediate future. As a second open topic, it would be interesting to adapt existing model checking techniques (based on automata construction) for temporal logics to solve the problem of existence of temporal stable models. This was done for infinite traces in (Cabalar and Diéguez, 2011; Cabalar and Demri, 2011), but no similar method has been implemented for finite traces on TEL_f or DEL_f yet. The importance of having an efficient implementation of such a method is that it would allow deciding non-existence of a plan in a given planning problem, something not possible by current incremental solving techniques. Another interesting topic is the optimization of grounding in temporal ASP specifications as those handled by telingo. The current grounding of telingo is inherited from incremental solving in clingo and does not exploit the semantics of temporal expressions that are available now in the input language. Finally, we envisage to extend the telingo system with features of DEL in order to obtain a powerful system for representing and reasoning about dynamic domains, not only providing an effective implementation of TEL and DEL but, furthermore, a platform for action and control languages, like A, B, C (Gelfond and Lifschitz, 1998; Giunchiglia et al., 2004) or GOLOG (Levesque et al., 1997).

References
Levesque, H.; Reiter, R.; Lespérance, Y.; Lin, F.; and Scherl, R. 1997. GOLOG: A logic programming language for

