Semantic Web

Part 4

The ontology layer 2: Reasoning in OWL

Riccardo Rosati

Corso di laurea magistrale in Ingegneria Informatica Sapienza Università di Roma 2012/2013

The Semantic Web Tower



OWL language

- Three species of OWL
 - OWL full is union of OWL syntax and RDF
 - OWL DL restricted to FOL fragment
 - OWL Lite is "easier to implement" subset of OWL DL
- OWL DL based on **SHIQ** Description Logic
 - In fact it is equivalent to $\frac{\text{SHOIN}(D_n)}{\text{DL}}$
- OWL DL Benefits from many years of DL research
 - Well defined semantics
 - Formal properties well understood (complexity, decidability)
 - Known reasoning algorithms
 - Implemented systems (highly optimised)

OWL class constructors

Constructor	DL Syntax	Example	Modal Syntax
intersectionOf	$C_1 \sqcap \ldots \sqcap C_n$	Human ⊓ Male	$C_1 \wedge \ldots \wedge C_n$
unionOf	$C_1 \sqcup \ldots \sqcup C_n$	Doctor ⊔ Lawyer	$C_1 \vee \ldots \vee C_n$
complementOf	$\neg C$	¬Male	$\neg C$
oneOf	$\{x_1\}\sqcup\ldots\sqcup\{x_n\}$	{john} ⊔ {mary}	$x_1 \lor \ldots \lor x_n$
allValuesFrom	$\forall P.C$	∀hasChild.Doctor	[P]C
someValuesFrom	$\exists P.C$	∃hasChild.Lawyer	$\langle P \rangle C$
maxCardinality	$\leqslant nP$	≤1hasChild	$[P]_{n+1}$
minCardinality	$\geqslant nP$	≥2hasChild	$\langle P \rangle_n$

Arbitrarily complex nesting of constructors:

– E.g., Person $\sqcap \forall$ hasChild.Doctor $\sqcup \exists$ hasChild.Doctor

DL knowledge bases (ontologies)

- An OWL ontology maps to a DL Knowledge Base $\mathcal{K} = \langle \mathcal{T}, \mathcal{A} \rangle$
 - $\mathcal{T}(TBox)$ is a set of axioms of the form:
 - $C \sqsubseteq D$ (concept inclusion)
 - $C \equiv D$ (concept equivalence)
 - $R \sqsubseteq S$ (role inclusion)
 - $R \equiv S$ (role equivalence)
 - $R^+ \sqsubseteq R$ (role transitivity)
 - \mathcal{A} (ABox) is a set of axioms of the form
 - $x \in D$ (concept instantiation)
 - $\langle x,y \rangle \in R$ (role instantiation)

DL vs. First-Order Logic

- in general, DLs correspond to decidable subclasses of firstorder logic (FOL)
- DL KB = first-order theory
- OWL Full is NOT a FOL fragment!
 - reasoning in OWL Full is undecidable
- OWL-DL and OWL-Lite are decidable fragments of FOL

DL vs. First-Order Logic

let $\mathcal{K} = \langle \mathcal{T}, \mathcal{A} \rangle$ be an ontology about persons where:

- *T* contains the following inclusion assertions: MALE □ PERSON FEMALE □ PERSON MALE □¬ FEMALE PERSON □∃Father¬.MALE
- *A* contains the following instance assertions:

MALE(Bob) PERSON (Mary) PERSON(Paul)

DL vs. First-Order Logic

• \mathcal{T} corresponds to the following FOL sentences: $\forall x. MALE(x) \rightarrow PERSON(x)$ $\forall x. FEMALE(x) \rightarrow PERSON(x)$ $\forall x. MALE(x) \rightarrow \neg FEMALE(x)$ $\forall x. PERSON(x) \rightarrow \exists y. Father(y,x) and MALE(y)$ \mathcal{A} corresponds to the following FOL ground atoms: MALE(Bob) **PERSON** (Mary) **PERSON(Paul)**

Inference tasks

- OWL-DL ontology = first-order logical theory
- verifying the formal properties of the ontology corresponds to **reasoning** over a first-order theory
- main reasoning tasks over ontologies:
 - consistency of the ontology
 - concept (and role) consistency
 - concept (and role) subsumption
 - instance checking
 - instance retrieval
 - query answering

Consistency of the ontology

- Is the ontology K=(T,A) consistent (non-selfcontradictory)?
- i.e., is there at least a model for K?
- intensional + extensional reasoning task
- fundamental formal property:
- inconsistent ontology => there is a semantic problem in K!
- K must be repaired

Consistency of the ontology

Example TBox:

```
MALE \sqsubseteq PERSON
FEMALE \sqsubseteq PERSON
MALE \sqsubseteq FEMALE
PERSON \sqsubseteq HasFather.MALE
PERSON \sqsubseteq HasMother.FEMALE
hasMother \sqsubseteq hasParent
hasFather \sqsubseteq hasParent
HasFather \sqsubseteq hasParent
```

Consistency of the ontology

- Example ABox:
 - MALE(Bob) MALE(Paul) FEMALE(Ann) hasFather(Ann,Paul) hasMother(Paul,Mary) BLACK-EYES(Mary)
 - \neg BLACK-EYES(Ann)
- ⇒ TBox + ABox **inconsistent** (Ann should have black eyes)

Concept consistency

- is a concept definition C consistent in a TBox T?
- i.e., is there a model of T in which C has a nonempty extension?
- intensional (schema) reasoning task
- detects a fundamental modeling problem in T:
 - if a concept is not consistent, then it can never be populated!

Concept subsumption

- is a concept C subsumed by another concept D in T?
- i.e., is the extension of C contained in the extension of D in every model of T?
- intensional (schema) reasoning task
- allows to do classification of concepts (i.e., to construct the concept ISA hierarchy)

Instance checking

- is an individual a a member of concept C in K?
- i.e., is the fact C(a) satisfied by every interpretation of K?
- intensional + extensional reasoning task
- basic "instance-level query" (tell me if object a is in class C)

Instance retrieval

- find all members of concept C in K
- i.e., compute all individuals a such that C(a) is satisfied by every interpretation of K
- intensional + extensional reasoning task
- (slight) generalization of instance checking

Query answering

- compute the answers to a **query** q in K (expressed in some query language)
- i.e., compute all tuples of individuals t such that q(t) is entailed by K (= q(t) is satisfied by every interpretation of K)
- extensional + extensional reasoning task
- generalization of instance checking and instance retrieval
- e.g.: database queries (SQL-like) over ontologies (or SPARQL-like queries)

Conjunctive queries

classes of queries over DL ontologies considered:

- **conjunctive queries** = subclass of SQL queries
 - correspond to select-project-join queries
- unions of conjunctive queries
 - correspond to select-project-join-union queries
- SPARQL-like queries
 - restrictions/extensions of SPARQL
- more expressive queries (e.g., SPARSQL)

Computational aspects of reasoning

- reasoning in OWL-DL is decidable (and the complexity is characterized)
- however: high computational complexity (EXPTIME)
- (optimized) reasoning algorithms developed
- OWL-DL reasoning tools implemented

Current OWL technology

two kinds of tools:

- OWL editors ("environments")
- OWL reasoners

OWL editors

- allow for visualizing/browsing/editing OWL ontologies
- able to connect to an external OWL reasoner
 => OWL "environments"
- main current tools:
 - Protege
 - SWOOP
 - OWLed2

OWL reasoning tools

two categories:

- OWL-DL reasoners
 - Racer, RacerPro
 - Pellet
 - Fact++
 - KAON2
- reasoners for "tractable fragments" of OWL-DL
 - QuOnto
 - OntoSearch2

OWL-DL reasoning tools

- all tools support "standard" reasoning tasks, i.e.:
 - consistency of the ontology
 - concept consistency
 - concept subsumption and classification
 - instance checking and retrieval
- they do not fully support conjunctive queries
- problem: the "official" query language for OWL has not been defined yet

Limits of current OWL-DL reasoners

- performance of OWL-DL reasoners:
- "practically good" for the intensional level
 - the size of a TBox is not likely to scale up too much
- not good for the extensional level
 - unable to handle instances (ABoxes) of large size (or even medium size)...
 - ...even for the basic extensional service (instance checking)

Limits of current OWL-DL reasoners

- why are these tools so bad with (large) ABoxes?
- two main reasons:
- current algorithms are mainly derived by algorithms defined for purely intensional tasks
 - no real optimization for ABox services
- these algorithms work in main memory
 => bottleneck for very large instances

OWL-DL technology vs. large instances

- the current limits of OWL-DL reasoners make it impossible to use these tools for real data integration on the web
- web sources are likely to be data intensive sources
- e.g., relational databases accessed through a web interface
- on the other hand, data integration is the prominent (future) application for Semantic Web technology!
 [Berners-Lee et al., IEEE Intelligent Systems, May 2006]

A solution: tractable OWL fragments

- how to overcome these limitations if we want to build data-intensive Semantic Web applications?
- solution 1 : do not reason over ontologies
- solution 2: limit the expressive power of the ontology language

=> tractable fragments of OWL

- solution 3: wait for more efficient OWL-DL reasoners
- to arrive at solution 2, we may benefit from the new technology developed for OWL tractable fragments

Tractable OWL fragments

- idea: sacrifice part of the expressiveness of the ontology language...
- ...to have more efficient ontology tools
- OWL Lite is a standardized fragment of OWL-DL
- is OWL Lite OK?
- NO! it is still too expressive for ABox reasoning
- OWL Lite is not really "lite"!

Tractable OWL fragments

- other fragments of OWL-DL have been proposed
- open problem (no standard yet)
- main current proposals:

OWL 2 PROFILES

- OWL 2 QL based on the DL DL-Lite
- OWL 2 EL based on the DL EL
- OWL 2 RL based on the DL DLP

References

- OWL W3C Web site: http://www.w3.org/2004/OWL/
- OWL overview:

http://www.w3.org/TR/owl-features/

- OWL resoning examples: http://owl.man.ac.uk/2003/why/latest/
- OWL Working Group (OWL 2): http://www.w3.org/2007/OWL/wiki/OWL_Working_Group
- OWL 2 profiles:

http://www.w3.org/TR/owl2-profiles/

• Web page on Description Logic reasoners: http://www.cs.man.ac.uk/~sattler/reasoners.html

References

• Racer (OWL reasoning tool):

http://www.racer-systems.com/

- Pellet (OWL reasoning tool): http://clarkparsia.com/pellet/
- KAON2 (OWL reasoning tool): http://kaon2.semanticweb.org/
- Protege (OWL ontology editor): http://protege.stanford.edu/
- SWOOP (OWL ontology editor):

http://code.google.com/p/swoop/